Formal Verification of Security Protocols with ProVerif

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Outline

Introduction

ProVerif

Protocol Formal Verification - Study Case

Conclusion

Questions

Cryptographic Protocols:

- series of steps;
- message exchanges;
- hostile environment;
- security properties.

Security Properties:

- secrecy;
- authenticity;
- \bullet integrity;
- ..

Security mechanisms that are commonly used by encryption protocols:

- Public Key Encryption
- Symmetric Encryption
- Hash Functions

The effectiveness of the protocol relies on keeping in secret the

The effectiveness of the protocol relies on keeping in secret the **keys**, not the **steps**.

Introduction Formal Methods

Formal methods are **techniques** used to model complex systems as **mathematical** and **logical** entities.

Introduction

Formal Methods Applied to Cryptographic Protocols

Aim: identify possible vulnerabilities!

Introduction

Formal Methods Applied to Cryptographic Protocols

Security Protocols:

- simple execution flow;
- difficult to design non-exploitable steps;

Introduction

Formal Methods Applied to Cryptographic Protocols

Initial research on security protocols formal verification date back from the 80's. Nowadays there are multiple automated formal verification tools.

ProVerif Introduction to the Tool

ProVerif is based on the formal model (Dolev-Yao model).

Introduction to the Tool - Dolev-Yao model

Attackers are capable of:

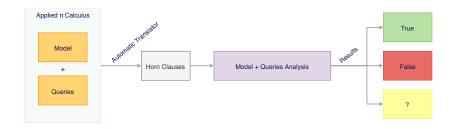
- permeating themselves in **between the communication** of two participants in any process of the protocol;
- modifying and copying fragments of information sent in the network;
- replicating messages;
- forging messages;

Introduction to the Tool - Dolev-Yao model

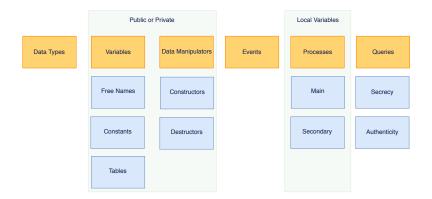
Attackers are capable of:

- keeping track of all messages sent in the network;
- actively participating as normal agents in the protocol;
- receiving responses sent to other participants

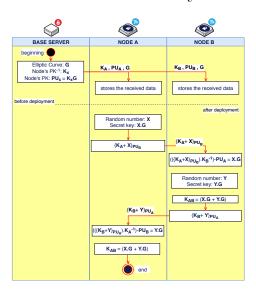
Introduction to the Tool



Code Organization



Saqib, 2016



Saqib, 2016 - Model Breakdown

Data Types

```
type skey.
type secretkey.
type id.
type pkey.
```



```
free C1: channel.
free C2: channel.
free CX: channel.
const GP: g [private].
table alreadyPaired(id, id).
```



```
let nodeA() =
                               new KA: skev: new IDA: id:
                               let PUA = calcPublicKey(KA, GP) in
                               insert publicKeysA(IDA, PUA); out(CX, PUA);
            Local Variables
                           5
                               !(
                                   get publicKeysB(IDBR1: id, PURB: pkey) in (
                                   get alreadyPaired(=IDA, =IDBR1) in (event eNodesAlreadyPairedA(
Events
             Processes
                                          → IDA, IDBR1))
                                   else (
                                   new X: number:
               Main
                          10
                                   let XG = calcSecretKev(X, GP) in
                          11
                                   event eNodeACreatesTheSecretKev(XG):
                          12
                                   let KAXPUB = encryption( addsSkeyPlusNumber(KA, X), PURB) in
                          13
                                   event eSendKAXPUB(KAXPUB);
             Secondary
                          14
                                   out(C1, (KAXPUB, IDA, IDBR1));
                          15
                                   in(C2, (K2:g, IDBR2: id, IDAR: id));
                                   if ((TDAR = TDA) && (TDBR2 = TDBR1)) then
                          16
                          17
                                   let GY = decryption(K2, KA, PURB) in(
                          18
                                   event eDecryptedUsingKA(GY,KA):
                          19
                                   let XGGY = calcSharedKey(XG, GY) in
                                   insert alreadyPaired(IDA, IDBR1):
                          20
                          21
                                   event eNodeAComputesSharedKev(XGGY))))
                          22
```



```
1 process
2 (!nodeB() | !nodeA())
```

Saqib, 2016 - Secrecy Queries

"Unauthorized agents are not capable of deriving specific information..."

Saqib, 2016 - Secrecy Queries



Secrecy

```
1 query secret KAXPUB.
2 query secret KBYPUA.
3 query attacker(new KA).
```

4 query attacker(new KB).

Authenticity

Protocol Formal Verification - Study Case Saqib, 2016 - Authenticity Queries

Correspondence Assertions "if an event **e** has been executed, then **e**' has been previously executed"

Protocol Formal Verification - Study Case Saqib, 2016 - Authenticity

Aliveness

Weak Agreement Non-injective Agreement Injective Agreement

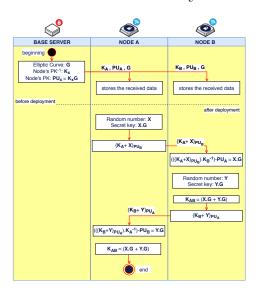
Saqib, 2016 - Authenticity Queries



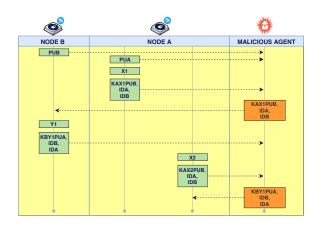
Protocol Formal Verification - Study Case Sagib, 2016 - Model Breakdown

Did the **symmetric key** calculated by **A** contain the secret key calculated by **B**, and did the **symmetric key** calculated by **B** actually contain the secret key sent by **A**?

Saqib, 2016



Saqib, 2016 - Result Analysis



Conclusion "Take Home" Messages

- It is **difficult** to design security protocols with **no vulnerabilities**;
- Formal verification techniques help on the process of checking if protocols guarantee certain security properties;
- ProVerif is one of the available tools to automate formal verification;

Conclusion "Take Home" Messages

- ProVerif can only verify what the user provides to it.
- Find the perfect balance between the levels of detail and abstraction for your model;
- The results given by ProVerif should serve as a **tool** to improve the analysis of the protocol.

Questions?